

Hybrid Best First Search for Computing “Anytime” Partition Function

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Definition

- 1 Set $X = \{X_1, \dots, X_n\}$ of **variables**, with a **domain** D_i containing values (boolean, integer, real).
- 2 Set Φ of **local functions** ϕ_S involving variables of $S \subset X$ (scope).
- 3 A **joint function** for a full assignment t :

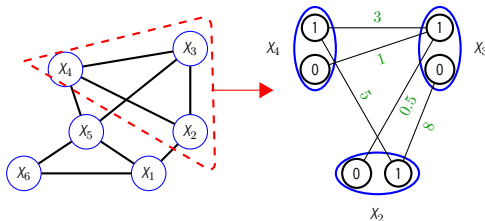
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Toy Example



Probability of an assignment

The joint **probability** of a **complete assignment** is defined as:

$$p(t) = \frac{P(t)}{Z} = \frac{1}{Z} \prod_{\phi_S \in \Phi} \phi_S(t[S])$$

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Partition Function

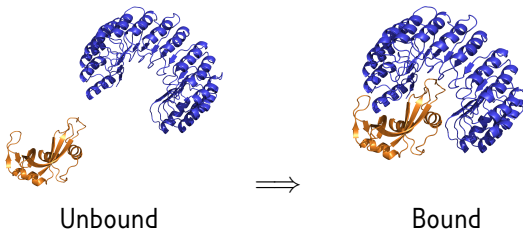
The **normalizing constant** or **Partition Function**:

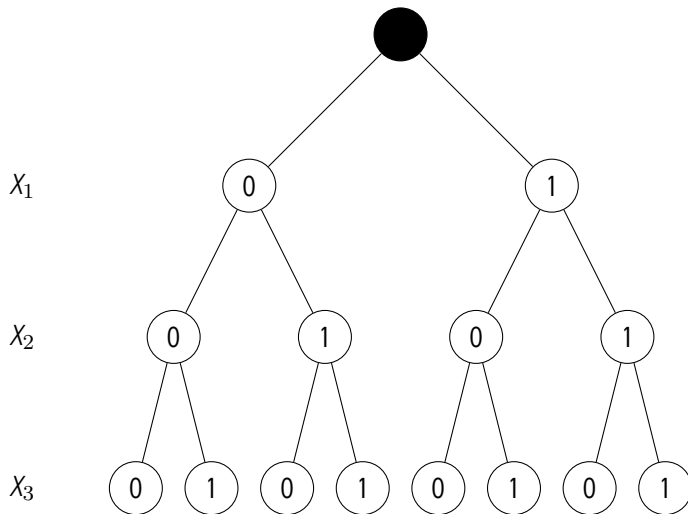
$$Z = \sum_t P(t) = \sum_t \prod_{\phi_S \in \Phi} \phi_S(t[S])$$

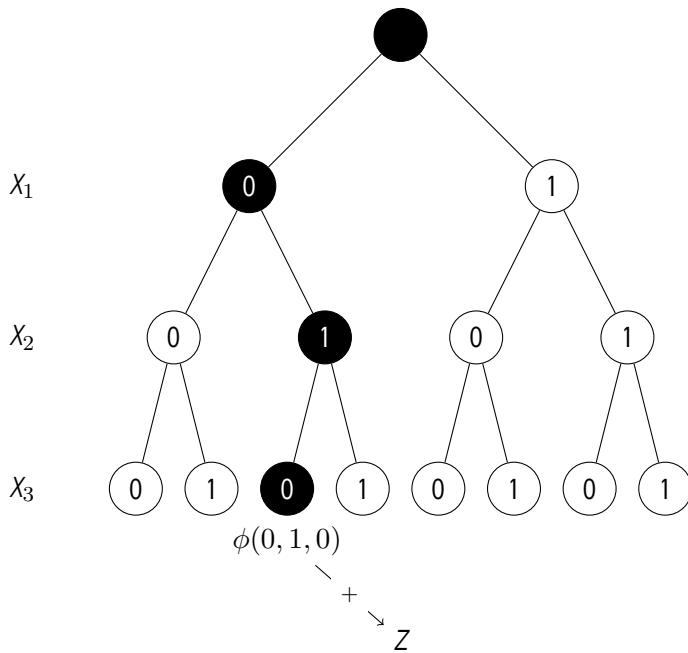
Constante d'affinité

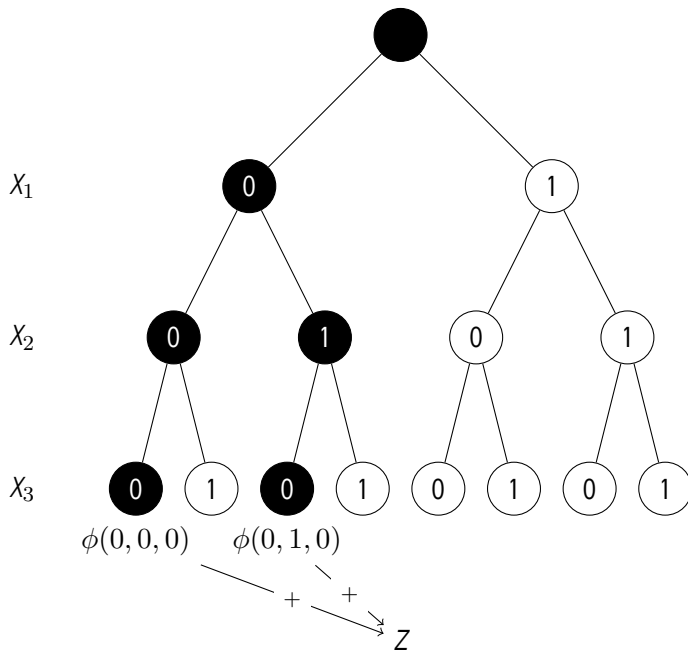
We can **approximate** the affinity between two proteins P_1 et P_2 forming a complex C by:

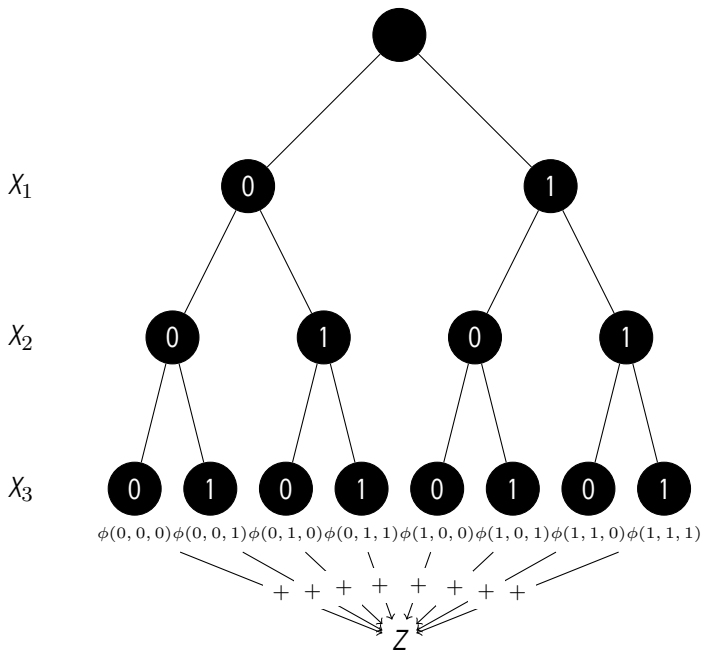
$$K_a = e^{(k_B T)} \frac{C^0}{8\pi^2} \frac{\sigma_{P_1} \sigma_{P_2}}{\sigma_C} \frac{Z(T, V, C)}{Z(T, V, P_1) Z(T, V, P_2)}$$







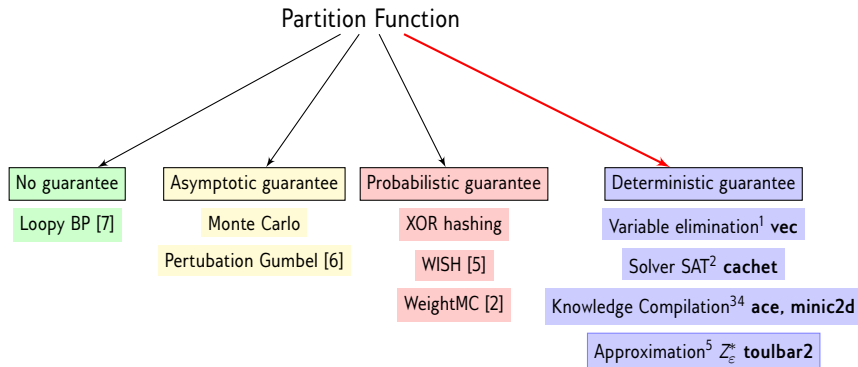




Combinatorial Explosion

Computing Z is classified as **#P-complete** problem so it is really hard





¹Rina Dechter. "Bucket Elimination: A Unifying Framework for Reasoning". In: *Artificial Intelligence* 113.1-2 (1999), pp. 41-85.

²Tian Sang, Paul Beame, and Henry A Kautz. "Performing Bayesian inference by weighted model counting". In: *AAAI*. vol. 5. 2005, pp. 475-481.

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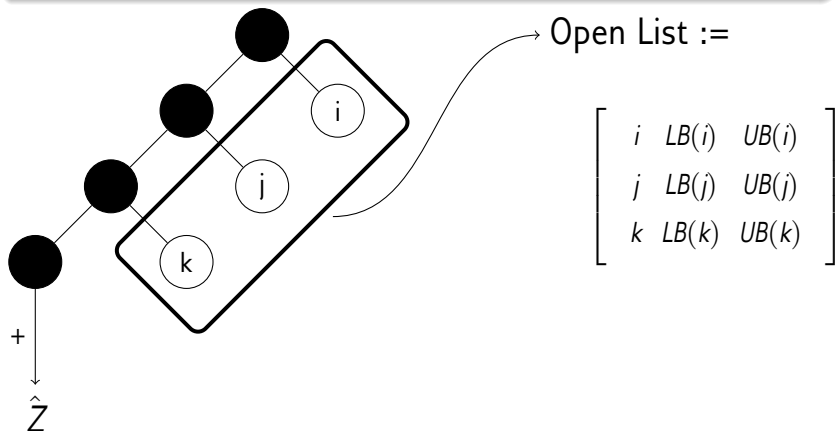
⁴Umut Oztok and Adnan Darwiche. "A top-down compiler for sentential decision diagrams". In: *Proceedings of the 24th International Conference on Artificial Intelligence*. AAAI Press. 2015.

⁵Clément Viricel et al. "Guaranteed weighted counting for affinity computation: Beyond determinism and structure". In: *International Conference on Principles and Practice of Constraint Programming*. Springer. 2016, pp. 733-750.

Hybrid Best First Search

We adapt **HBFS^a** search algorithm to **compute the partition function**

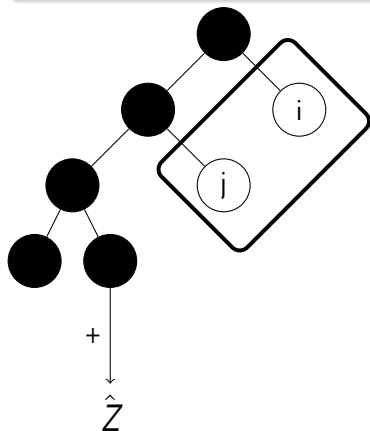
^aDavid Allouche et al. "Anytime hybrid best-first search with tree decomposition for weighted CSP". . In: *International Conference on Principles and Practice of Constraint Programming*. Springer. 2015, pp. 12–29.



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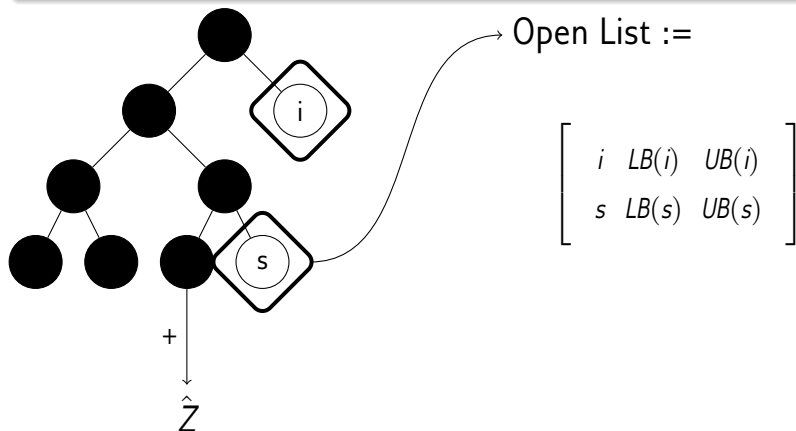
Open List :=

$$\begin{bmatrix} i & LB(i) & UB(i) \\ j & LB(j) & UB(j) \end{bmatrix}$$

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Function HBFS-C

```
Open = [(root, LB(root), UB(root))] ;
```

```
while Open  $\neq \emptyset$  do
```

```
    n = pop(Open) ;
```

```
    restore(n) ;
```

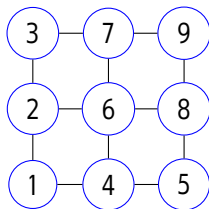
```
     $\hat{Z} = DFS(n, \hat{Z}, k)$  ;
```

$$\hat{Z} + \sum_{n \in Open} LB(n) \leq Z \leq \hat{Z} + \sum_{n \in Open} UB(n)$$

We can have an anytime guarantee with:

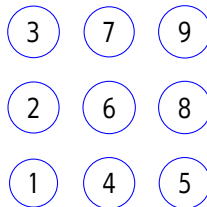
$$\hat{Z} + \sum_{n \in Open} UB(n) \leq (1 + \varepsilon) \left(\hat{Z} + \sum_{n \in Open} LB(n) \right)$$

Original graph



$$P \propto \prod_t \phi_t(x_t)$$

Naive Mean Field



$$Q \propto \prod_{i \in V} q_i(x_i)$$

$$\log(Z) \geq \sum_{i \in V} \sum_{x_i} [S_i(q_i) - q_i(x_i) E_i(x_i)] - \sum_{(i,j) \in E} \sum_{x_i, x_j} q_i(x_i) q_j(x_j) E_{ij}(x_i, x_j)$$

Function MF-LB(n)

$t \leftarrow 0$;

Initialise $q^{(t)}$;

while $q^{(t)}$ converge **do**

for $i \in X(n)$ **do**

$q_i^{(t+1)} \leftarrow \exp \left(-E_i(x_i) - \sum_{j \in \mathcal{N}(i)} \sum_{x_j} q_j(x_j) E_{ij}(x_i, x_j) \right)$;

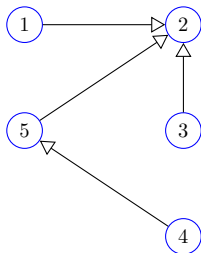
$Z_i \leftarrow \sum_{x_i} q_i^{t+1}(x_i)$;

$q_i^{(t+1)} \leftarrow \frac{q_i^{(t+1)}}{Z_i}$;

Upper Bound by Maximum Spanning Tree

Define a **maximum spanning tree** $T \subset \Phi$ and by applying **dynamic programming** to $T' = T \cup \{E_S \in \Phi : |S| < 2\}$, we have an exact $Z_{T'}$ in **polynomial time**.

$$Z \leq Ub_T = \underbrace{\left(\sum_{t \in D^X} \prod_{\phi_S \in T} \phi_S(t) \right)}_{\text{programmation dynamique}} \cdot \left(\prod_{\phi_S \in \Phi \setminus T} \max_{t \in D^S} \phi_S(t) \right)$$



Extension

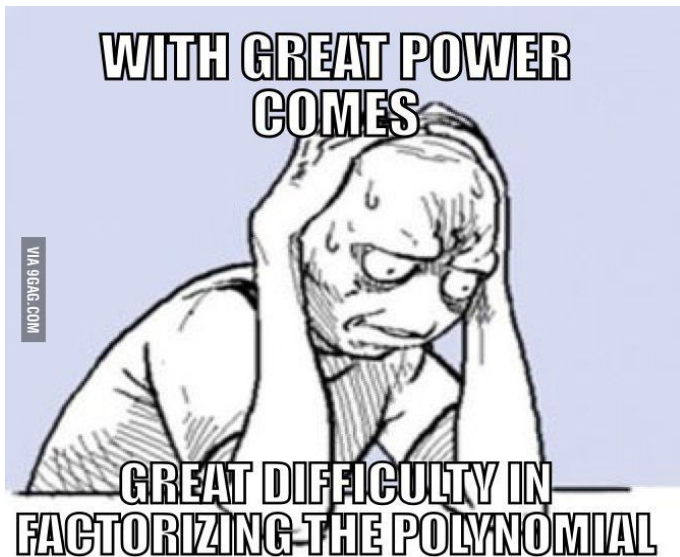
- Mixing BTD (Backtrack Tree decomposition) with HBFS-C
- Integrate Z_{ϵ}^* pruning to HBFS

Test

Run a battery of tests (HBFS-C; HBFS-C + Z_{ϵ}^* ; HBFS-C+BTB; HBFS-C + BTB+ Z_{ϵ}^*) to see the dynamic of all the algorithm.

Application

If these are improvement then try to solve large protein instances to predict affinity.





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