# Modèles Probabilistes Complexes 

Models et Algorithms

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## Inference in Bayesian networks

## Variable eliminitation

Input : a set of CPTs $\mathbf{P}$ and a set of variables $\mathbf{X}$
Output: $P(\mathbf{X})$
$\mathbf{1} \mathbf{W} \leftarrow$ all the variables of the CPTs of $\mathbf{P}$ except $\mathbf{X}$
2 while $W \neq \emptyset$ do
3 let $X_{j}$ be a variable in $\mathbf{W}$; remove $X_{j}$ from $\mathbf{W}$
$4 \quad$ let $\mathbf{Q}$ be the set of tables in $\mathbf{P}$ containing $X_{j}$
$5 \quad$ compute table $q=\sum_{x_{j}} \prod_{f \in \mathbf{Q}} f$
6
$\mathbf{P} \leftarrow(\mathbf{P} \backslash \mathbf{Q}) \cup\{q\}$
7 return table $\prod_{f \in \mathbf{P}} f$

## Dynamic Bayesian networks (dBNs)

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$$
\text { 2-TBN } \quad \text { dBN of length } 2
$$



Moralization


## Problems to fix

(1) The dBN can be arbitrarily large,
(2) Need to re-triangulate the dBN if its length changes,
(3) The size of some cliques may be a function of $T$.

## Constrained triangulation : forward

How to provide a way to avoid re-triangulating the dBN ?

## Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

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$\hookrightarrow$ Backward interface: nodes with parents or spouses in the previous slice.

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$\hookrightarrow$ Forward interface: nodes with children in the next slice.

## General interface




Slice $T$

- Transition model :

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## Interface

Subset of nodes such that if they were removed they will disconnect the past from the future in the moral graph of a dBN.

## Complexity for interface based ceo

## Interface based Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order given by an interface.
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Best case : clique of size $|I|+1$

## Upper bound



Worst case : clique of size $|I|+S$ ( $S$ : \# nodes in a slice)

## How to improve the bounds?


$\omega=2$

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## Relation to the Minimum s-t CuT problem

## Theorem

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$\Rightarrow$ Polynomial time algorithm to find the minimum interface in a dBN of length $T$

## Building the optimal elimination order

For wich value of $T$ do we have the minimum interface?

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Iteration 1 Iteration 2 Iteration 3
Optimal solution


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## Theorem

At most $O(h)$ iterations where $h$ is the number of nodes in a slice.

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## Better lower bound

Best case : clique of size $\left|I^{*}\right|+1$

## Better upper bound

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Maximum clique size :


| Fig. | B-SS | F-SS | Min-Elim |
| :--- | :---: | :---: | :---: |
| C | 3.16 | 3.03 | 3.03 |
| B | 5.54 | 5.54 | 5.54 |
| D | 1.38 | 3.00 | 1.38 |
| A | 2.07 | 3.45 | 2.07 |
| E | 3.46 | 3.23 | 3.00 |

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Mean clique size :

| Fig. | B-SS | F-SS | Min-Elim |
| :--- | :---: | :---: | :---: |
| C | 3.46 | 3.46 | 3.46 |
| B | 5.54 | 5.54 | 5.54 |
| D | 1.39 | 3.46 | 1.39 |
| A | 2.08 | 3.46 | 2.08 |
| E | 4.16 | 3.46 | 3.46 |

## Experimental results - Random dBNs

Results on randomly generated dBNs
(1) 5 variables per slices
(2) Cardinalities chosen uniformly between 2 and 8
(3) DBNs unrolled 500 time steps

Mean clique size
Interface size


## Experimental results - Random dBNs

Results on randomly generated dBNs
(1) 10 variables per slices
(2) Cardinalities chosen uniformly between 2 and 8
(3) DBNs unrolled 500 time steps

Mean clique size


Interface size


## Experimental results - Random dBNs

Results on randomly generated dBNs
(1) 15 variables per slices
(2) Cardinalities chosen uniformly between 2 and 8
(3) DBNs unrolled 500 time steps


Interface size


## Motivation



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```
our goal : design complex PGM speed-up inference in PRMs
```


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PRMs are an extension of object-oriented Bayesian Networks (Pfeffer \& Koller, 1997) and (Bangsø \& Wuillemin, 2000).


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- Reference between classes are used to define dependencies between different fragment.
- Classes are instantiated in a system.



## (1) Reinforcing the object-oriented aspect of PRMs

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## (2) Interface and multiple inheritance

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## (2) Interface and multiple inheritance



## System : the printer example

## A Class Dependency Graph.



- Complex reference


## A Relational Skeleton.



## Applications

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- All of which have been successfully used by experts to represent different complex systems.
- The models we created required the modification of the PRM framework presented here.


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- Faster inference!


## Overview of PRMs



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Relational skeleton

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- However, it works only if the inner attributes are not observed.

Once inner attributes are eliminated, SVE proceeds with a bottom-up elimination of each instance.

## Structured Inference

Probability $P(C)$ ?


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Elimination at class level : performed only once!


It works only if the inner attributes are not observed.
Only internal variables can be eliminated at class level!

## D-Separation in PRMs ([oti and Werilemin 2010)



Figure: Different configurations due to evidence.

## Integrating D-Separation in SVE

Following SVE bottom-up elimination order :
(1) Recursive calls are made on each reverse slot chain.
(2) If at least one reverse slot chain is active :
(1) Activate any required slot chains.
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Generalized BayesBall rule

## Experimental results

## The example used for our experimentations

- We used the printer example with the following parameters:
- different number of computers per room;
- different number of printers per room ;
- different number of evidences.
- The examples we used have a 1400 to 24000 attributes.


## Efficiency of D-Separation



## Increasing the number of computers



## Increasing the number of printers



## Only internal variables can be eliminated at class level

- Construction by experts $\Longrightarrow$ genericity and design patterns.


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## Substitution rule



## Rule

One random variable can only belong to one instance.

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## Optimal dynamic classes

## Proposition

The following problem is NP-hard :
Instance : a PRM, an integer $K \geq 0$
Question : does there exist a set of dynamic classes / substitutions s.t. the number of operations performed by structured inference is $\leq K$ ?

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Instance : a PRM, an integer $K \geq 0$
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remains NP-hard even when classical inference is polynomial

## Approximate algorithm

## Boundary graph



Relational skeleton

## Approximate algorithm

## Boundary graph


$\Longrightarrow$ find frequent patterns in the boundary graph

## Search tree

- Mining frequent subgraphs

> [Inokuchi et al, 05], [Kuramochi and Karypis, 01], [Yan and Han, 02]

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- Mining frequent subgraphs
[Inokuchi et al, 05], [Kuramochi and Karypis, 01], [Yan and Han, 02]
- Variant of the gSpan algorithm :

1 edge
2 edges
3 edges
4 edges


- Node $=$ (dynamic class, set of instances)


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- Node $=($ dynamic class, set of instances $)$
- Possible substitutions (Rule 1) : Max Independence Set


## Another pruning rule

## Efficiency Pruning Rule

Prune nodes/classes whose substitutions do not speed-up inference.

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Prune nodes/classes whose substitutions do not speed-up inference.

- Gain estimation by dynamic programming
$\Longrightarrow \alpha$-value : $\alpha>0 \Longleftrightarrow$ class unattractive
search tree not monotonically $\alpha$-decreasing!
- Rule applied : prune subtree whenever $\alpha>0$


## Experiments



