## Modèles Probabilistes Complexes Models et Algorithms

## Pierre-Henri WUILLEMIN

Christophe Gonzales - Lionel Torti - Morgan Chopin

(prenom.nom@lip6.fr)

#### 8 novembre 2011







### Variable eliminitation

```
Input : a set of CPTs \mathbf{P} and a set of variables \mathbf{X}
Output : P(\mathbf{X})
```

 $\mathbf{1} \ \mathbf{W} \leftarrow \mathsf{all} \ \mathsf{the} \ \mathsf{variables} \ \mathsf{of} \ \mathsf{the} \ \mathsf{CPTs} \ \mathsf{of} \ \mathbf{P} \ \mathsf{except} \ \mathbf{X}$ 

2 while 
$$\mathbf{W} \neq \emptyset$$
 do

3 let  $X_j$  be a variable in  $\mathbf{W}$ ; remove  $X_j$  from  $\mathbf{W}$ 

4 let 
$$\mathbf{Q}$$
 be the set of tables in  $\mathbf{P}$  containing  $X_j$ 

compute table 
$$q = \sum_{X_j} \prod_{f \in \mathbf{Q}} f$$

6 
$$\mathbf{P} \leftarrow (\mathbf{P} \backslash \mathbf{Q}) \cup \{q\}$$

7 return table  $\prod_{f \in \mathbf{P}} f$ 





- Transition model :  $p(\mathbf{X}_t | \mathbf{X}_{t-1}) = p(\mathbf{X}_1 | \mathbf{X}_0)$   $= \prod_{i=0}^{n-1} p(X_1^i | pa(X_1^i))$
- Prior distribution :  $p(\mathbf{X}_0) = \prod_{i=0}^{n-1} p(X_0^i | pa(X_0^i))$

Bayesian network that models a dynamic system (or a repetition of patterns) :



• Transition model :  $p(\mathbf{X}_t | \mathbf{X}_{t-1}) = p(\mathbf{X}_1 | \mathbf{X}_0)$   $= \prod_{i=0}^{n-1} p(X_1^i | pa(X_1^i))$ 

• Prior distribution :  $p(\mathbf{X}_0) = \prod_{i=0}^{n-1} p(X_0^i | pa(X_0^i))$ 



- Transition model :  $p(\mathbf{X}_t | \mathbf{X}_{t-1}) = p(\mathbf{X}_1 | \mathbf{X}_0)$   $= \prod_{i=0}^{n-1} p(X_1^i | pa(X_1^i))$
- Prior distribution :  $p(\mathbf{X}_0) = \prod_{i=0}^{n-1} p(X_0^i| pa(X_0^i))$

Bayesian network proposes a simplification of the dynamic modelisation :





• Prior distribution :  

$$p(\mathbf{X}_0) = \prod_{i=0}^{n-1} p(X_0^i | pa(X_0^i))$$

Bayesian network proposes a simplification of the dynamic modelisation :



- Transition model :  $p(\mathbf{X}_t | \mathbf{X}_{t-1}) = p(\mathbf{X}_1 | \mathbf{X}_0)$   $= \prod_{i=0}^{n-1} p(X_1^i | pa(X_1^i))$
- Prior distribution :  $p(\mathbf{X}_0) = \prod_{i=0}^{n-1} p(X_0^i| pa(X_0^i))$

Bayesian network proposes a simplification of the dynamic modelisation :





• Prior distribution :  

$$p(\mathbf{X}_0) = \prod_{i=0}^{n-1} p(X_0^i | pa(X_0^i))$$

Bayesian network proposes a simplification of the dynamic modelisation :



• Transition model :  $p(\mathbf{X}_t | \mathbf{X}_{t-1}) = p(\mathbf{X}_1 | \mathbf{X}_0)$   $= \prod_{i=0}^{n-1} p(X_1^i | pa(X_1^i))$ 

• Prior distribution :  

$$p(\mathbf{X}_0) = \prod_{i=0}^{n-1} p(X_0^i | pa(X_0^i))$$

### How to triangulate a dBN?

A dBN is a Bayesian network  $\Rightarrow$  triangulation based inference algorithm.

### How to triangulate a dBN?

A dBN is a Bayesian network  $\Rightarrow$  triangulation based inference algorithm.

• Unroll the 2-TBN to any desired length T and moralize.

### How to triangulate a dBN?

A dBN is a Bayesian network  $\Rightarrow$  triangulation based inference algorithm.

- Unroll the 2-TBN to any desired length T and moralize.
- Apply any classical triangulation algorithm.

### How to triangulate a dBN?

A dBN is a Bayesian network  $\Rightarrow$  triangulation based inference algorithm.

- Unroll the 2-TBN to any desired length T and moralize.
- Apply any classical triangulation algorithm.



### How to triangulate a dBN?

A dBN is a Bayesian network  $\Rightarrow$  triangulation based inference algorithm.

- Unroll the 2-TBN to any desired length T and moralize.
- Apply any classical triangulation algorithm.



### Problems to fix

### How to triangulate a dBN?

A dBN is a Bayesian network  $\Rightarrow$  triangulation based inference algorithm.

- Unroll the 2-TBN to any desired length T and moralize.
- Apply any classical triangulation algorithm.



### Problems to fix

The dBN can be arbitrarily large,

### How to triangulate a dBN?

A dBN is a Bayesian network  $\Rightarrow$  triangulation based inference algorithm.

- Unroll the 2-TBN to any desired length T and moralize.
- Apply any classical triangulation algorithm.



#### Problems to fix

- The dBN can be arbitrarily large,
- Need to re-triangulate the dBN if its length changes,

## How to triangulate a dBN?

A dBN is a Bayesian network  $\Rightarrow$  triangulation based inference algorithm.

- Unroll the 2-TBN to any desired length T and moralize.
- Apply any classical triangulation algorithm.



#### Problems to fix

- The dBN can be arbitrarily large,
- Need to re-triangulate the dBN if its length changes,
- The size of some cliques may be a function of T.

How to provide a way to avoid re-triangulating the dBN?

#### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

How to provide a way to avoid re-triangulating the dBN?

#### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

### The forward slice-by-slice elimination (F-SS) [dereviene, 2001]

How to provide a way to avoid re-triangulating the dBN?

### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

### The forward slice-by-slice elimination (F-SS) [basedite, 2001



How to provide a way to avoid re-triangulating the dBN?

### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

### The forward slice-by-slice elimination (F-SS) [basedite, 2001



How to provide a way to avoid re-triangulating the dBN?

### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

### The forward slice-by-slice elimination (F-SS) [basedite, 2001



How to provide a way to avoid re-triangulating the dBN?

### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

### The forward slice-by-slice elimination (F-SS) [basedite, 2001



How to provide a way to avoid re-triangulating the dBN?

### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

### The forward slice-by-slice elimination (F-SS) [basedite, 2001



How to provide a way to avoid re-triangulating the dBN?

### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

### The forward slice-by-slice elimination (F-SS) [basedite, 2001



How to provide a way to avoid re-triangulating the dBN?

### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

### The forward slice-by-slice elimination (F-SS) [basedite, 2001



How to provide a way to avoid re-triangulating the dBN?

### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

### The forward slice-by-slice elimination (F-SS) [basedite, 2001



How to provide a way to avoid re-triangulating the dBN?

### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

### The forward slice-by-slice elimination (F-SS) (provided 2001

Consists in eliminating a slice after an other from the past to the future.



 $\hookrightarrow$  *Backward interface* : nodes with parents or spouses in the previous slice.

How to provide a way to avoid re-triangulating the  $\mathsf{dBN}\,?$ 

How to provide a way to avoid re-triangulating the dBN?

#### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

How to provide a way to avoid re-triangulating the dBN?

### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

## The backward slice-by-slice elimination (B-SS) plansing 2001



How to provide a way to avoid re-triangulating the dBN?

### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

## The backward slice-by-slice elimination (B-SS) plansing 2001



How to provide a way to avoid re-triangulating the dBN?

### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

## The backward slice-by-slice elimination (B-SS) plansing 2001



How to provide a way to avoid re-triangulating the dBN?

### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

## The backward slice-by-slice elimination (B-SS) provide 2001


## Constrained triangulation : backward

How to provide a way to avoid re-triangulating the dBN?

### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

### The backward slice-by-slice elimination (B-SS) provide 2001

Consists in eliminating a slice after an other from the future to the past :



## Constrained triangulation : backward

How to provide a way to avoid re-triangulating the dBN?

### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

### The backward slice-by-slice elimination (B-SS) provide 2001

Consists in eliminating a slice after an other from the future to the past :



# Constrained triangulation : backward

How to provide a way to avoid re-triangulating the dBN?

### Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order.

### The backward slice-by-slice elimination (B-SS) parameter 2000

Consists in eliminating a slice after an other from the future to the past :



 $\hookrightarrow$  Forward interface : nodes with children in the next slice.









#### Interface

Subset of nodes such that if they were removed they will disconnect the past from the future in the moral graph of a dBN.

### Interface based Constrained elimination order (ceo)

Elimination order such that we enforce nodes to be eliminated with respect to a partial order given by an interface.

Provide a way to avoid re-triangulating the dBN,

### Interface based Constrained elimination order (ceo)

- Provide a way to avoid re-triangulating the dBN,
- 2 The size  $\omega$  of the maximal clique is bounded :

### Interface based Constrained elimination order (ceo)

- Provide a way to avoid re-triangulating the dBN,
- 2 The size  $\omega$  of the maximal clique is bounded :



### Interface based Constrained elimination order (ceo)

- Provide a way to avoid re-triangulating the dBN,
- 2 The size  $\omega$  of the maximal clique is bounded :



### Interface based Constrained elimination order (ceo)

- Provide a way to avoid re-triangulating the dBN,
- 2 The size  $\omega$  of the maximal clique is bounded :



### Interface based Constrained elimination order (ceo)

- Provide a way to avoid re-triangulating the dBN,
- 2 The size  $\omega$  of the maximal clique is bounded :







 $\hookrightarrow$  Looking for a smaller interface by extending the slice [bilmes et al, 2003]



 $\hookrightarrow$  Looking for a smaller interface by extending the slice [bilmes et al, 2003]





 $\hookrightarrow$  Looking for a smaller interface by extending the slice [bilmes et al, 2003]

• Algorithm to find the optimal interface in polynomial time,



 $\hookrightarrow$  Looking for a smaller interface by extending the slice [bilmes et al, 2003]

Algorithm to find the optimal interface in polynomial time,
Deduce a *ceo* with better theoretical guarantee on cliques size.



 $\hookrightarrow$  Looking for a smaller interface by extending the slice [bilmes et al, 2003]

Algorithm to find the optimal interface in polynomial time,
Deduce a *ceo* with better theoretical guarantee on cliques size.

# Relation to the MINIMUM s-t CUT problem

#### Theorem

Given a dBN of length T, finding a minimum interface is equivalent to solve the minimum s-t cut problem

# Relation to the MINIMUM s-t CUT problem

#### Theorem

Given a dBN of length T, finding a minimum interface is equivalent to solve the minimum s-t cut problem



# Relation to the MINIMUM s-t CUT problem

#### Theorem

Given a dBN of length T, finding a minimum interface is equivalent to solve the minimum s-t cut problem



 $\Rightarrow$  Polynomial time algorithm to find the minimum interface in a dBN of length  $\mathcal{T}$  |

For wich value of T do we have the minimum interface?

For wich value of T do we have the minimum interface?  $\hookrightarrow$  Proceed by iteratively solving the min *s*-*t* cut problem and increasing the length T

For wich value of T do we have the minimum interface?  $\hookrightarrow$  Proceed by iteratively solving the min *s*-*t* cut problem and increasing the length T



For wich value of T do we have the minimum interface?  $\hookrightarrow$  Proceed by iteratively solving the min *s*-*t* cut problem and increasing the length T



#### Theorem

At most O(h) iterations where h is the number of nodes in a slice.

# Building the elimination order



# Building the elimination order



# Building the elimination order



#### Better lower bound

Best case : clique of size  $|I^*| + 1$ 

#### Better upper bound

Worst case : clique of size  $|I^*| + S$ (S : # nodes in a slice)

## Experimental results - Classical dBNs



## Experimental results - Classical dBNs



#### Maximum clique size :

Fig.	B-SS	F-SS	$\operatorname{Min-Elim}$
С	3.16	3.03	3.03
В	5.54	<b>5.54</b>	5.54
D	1.38	3.00	1.38
A	2.07	3.45	2.07
E	3.46	3.23	3.00

## Experimental results – Classical dBNs



#### Maximum clique size :

Fig.	B-SS	F-SS	$\operatorname{Min-Elim}$
С	3.16	3.03	3.03
В	5.54	<b>5.54</b>	5.54
D	1.38	3.00	1.38
A	2.07	3.45	2.07
E	3.46	3.23	3.00

#### Mean clique size :

Fig.	B-SS	F-SS	$\operatorname{Min-Elim}$
С	3.46	3.46	3.46
В	5.54	<b>5.54</b>	5.54
D	1.39	3.46	1.39
A	2.08	3.46	2.08
E	4.16	3.46	3.46

### Experimental results - Random dBNs

Results on randomly generated dBNs

- **0 5** variables per slices
- Cardinalities chosen uniformly between 2 and 8
- OBNs unrolled 500 time steps

Mean clique size





### Experimental results - Random dBNs

Results on randomly generated dBNs

- **1** variables per slices
- Cardinalities chosen uniformly between 2 and 8
- OBNs unrolled 500 time steps

Mean clique size







### Experimental results – Random dBNs

Results on randomly generated dBNs

- 15 variables per slices
- Cardinalities chosen uniformly between 2 and 8
- OBNs unrolled 500 time steps

Mean clique size







## Motivation








• Large-scale Bayesian networks "undesirable" features :



• Large-scale Bayesian networks "undesirable" features : • design cost



• Large-scale Bayesian networks "undesirable" features :

• design cost

maintenance cost



• Large-scale Bayesian networks "undesirable" features :

- design cost
- maintenance cost
- Inference times



[Mahoney & Laskey, 96], [Pfeffer et al., 99]

16 / 39

Models et Algorithms

• Large-scale Bayesian networks "undesirable" features :

- design cost
- maintenance cost
- inference times
- $\implies$  often inadequate for large scale applications

Modèles Probabilistes Complexes



[Mahoney & Laskey, 96], [Pfeffer et al., 99]

16 / 39

Models et Algorithms

• Large-scale Bayesian networks "undesirable" features :

- design cost
- maintenance cost
- inference times
- $\implies$  often inadequate for large scale applications

Modèles Probabilistes Complexes

• First-order logic extensions (e.g., MLN)

[Jaeger, 97], [Kersting & De Raedt, 01]

• First-order logic extensions (e.g., MLN)

[Jaeger, 97], [Kersting & De Raedt, 01]

• Entity-Relationship extensions (e.g., MEBN)

[Heckerman et al., 04], [Laskey, 08]

• First-order logic extensions (e.g., MLN)

```
[Jaeger, 97], [Kersting & De Raedt, 01]
```

• Entity-Relationship extensions (e.g., MEBN)

[Heckerman et al., 04], [Laskey, 08]

• Object-Oriented extensions (e.g., PRM)

[Koller & Pfeffer, 97], [Mahoney & Laskey, 96]

• First-order logic extensions (e.g., MLN)

```
[Jaeger, 97], [Kersting & De Raedt, 01]
```

• Entity-Relationship extensions (e.g., MEBN)

[Heckerman et al., 04], [Laskey, 08]

• Object-Oriented extensions (e.g., PRM)

[Koller & Pfeffer, 97], [Mahoney & Laskey, 96]

our goal : design complex PGM and speed-up inference in PRMs

PRMs are an extension of object-oriented Bayesian Networks (Pfeffer & Koller, 1997) and (Bangsø & Wuillemin, 2000).





PRMs are an extension of object-oriented Bayesian Networks (Pfeffer & Koller, 1997) and (Bangsø & Wuillemin, 2000).



PRMs are an extension of object-oriented Bayesian Networks (Pfeffer & Koller, 1997) and (Bangsø & Wuillemin, 2000).



- Classes are BNs fragments.
- Nodes are called attributes.



PRMs are an extension of object-oriented Bayesian Networks (Pfeffer & Koller, 1997) and (Bangsø & Wuillemin, 2000).

#### PRMs in 4 definitions

- Classes are BNs fragments.
- Nodes are called attributes.
- Reference between classes are used to define dependencies between different fragment.



PRMs are an extension of object-oriented Bayesian Networks (Pfeffer & Koller, 1997) and (Bangsø & Wuillemin, 2000).

#### PRMs in 4 definitions

- Classes are BNs fragments.
- Nodes are called attributes.
- Reference between classes are used to define dependencies between different fragment.

• Classes are instantiated in a system.



## (1) Reinforcing the object-oriented aspect of PRMs (more second variant, term or m)



# (1) Reinforcing the object-oriented aspect of PRMs (more second variants and main)



# (1) Reinforcing the object-oriented aspect of PRMs (more second variants and main)



# (1) Reinforcing the object-oriented aspect of PRMs (1000 accord volume (2000 area)



# (1) Reinforcing the object-oriented aspect of PRMs (non-second volume to a m)



# (1) Reinforcing the object-oriented aspect of PRMs (non-second volume to a m)



#### Interface implementation

• An interface is only defined by a set of attributes types and references.

- An interface is only defined by a set of attributes types and references.
- This implies the absence of a DAG or attribute's CPT in the interface definition.

- An interface is only defined by a set of attributes types and references.
- This implies the absence of a DAG or attribute's CPT in the interface definition.
- Any class implementing an interface guarantees the existence of the interface's attributes and references.

- An interface is only defined by a set of attributes types and references.
- This implies the absence of a DAG or attribute's CPT in the interface definition.
- Any class implementing an interface guarantees the existence of the interface's attributes and references.







## System : the printer example



• SKOOB is a consortium of industrials, research laboratories and experts which focus on the use of probabilistic graphical models in reliability and risk management.

- SKOOB is a consortium of industrials, research laboratories and experts which focus on the use of probabilistic graphical models in reliability and risk management.
- It has also lead to the definition of a declarative object oriented language for the specification of PRMs (the SKOOL language).

- SKOOB is a consortium of industrials, research laboratories and experts which focus on the use of probabilistic graphical models in reliability and risk management.
- It has also lead to the definition of a declarative object oriented language for the specification of PRMs (the SKOOL language).
- The enhancement presented here have been implemented in two different prototypes.

- SKOOB is a consortium of industrials, research laboratories and experts which focus on the use of probabilistic graphical models in reliability and risk management.
- It has also lead to the definition of a declarative object oriented language for the specification of PRMs (the SKOOL language).
- The enhancement presented here have been implemented in two different prototypes.
- All of which have been successfully used by experts to represent different complex systems.

- SKOOB is a consortium of industrials, research laboratories and experts which focus on the use of probabilistic graphical models in reliability and risk management.
- It has also lead to the definition of a declarative object oriented language for the specification of PRMs (the SKOOL language).
- The enhancement presented here have been implemented in two different prototypes.
- All of which have been successfully used by experts to represent different complex systems.
- The models we created required the modification of the PRM framework presented here.

## Inference in PRM
• Any system can be transformed in a Bayesian Network (called a grounded network).

- Any system can be transformed in a Bayesian Network (called a grounded network).
- The generation algorithm is polynomial [Pfeffer, 2000].

- Any system can be transformed in a Bayesian Network (called a grounded network).
- The generation algorithm is polynomial [Pfeffer, 2000].
- It is possible to use Bayesian Networks state-of-the-art inference algorithms.

- Any system can be transformed in a Bayesian Network (called a grounded network).
- The generation algorithm is polynomial [Pfeffer, 2000].
- It is possible to use Bayesian Networks state-of-the-art inference algorithms.

- Any system can be transformed in a Bayesian Network (called a grounded network).
- The generation algorithm is polynomial [Pfeffer, 2000].
- It is possible to use Bayesian Networks state-of-the-art inference algorithms.

- It is easy to produce very large scale systems :
  - naive grounded networks use too much memory;

- Any system can be transformed in a Bayesian Network (called a grounded network).
- The generation algorithm is polynomial [Pfeffer, 2000].
- It is possible to use Bayesian Networks state-of-the-art inference algorithms.

- It is easy to produce very large scale systems :
  - naive grounded networks use too much memory;
  - intelligent grounded networks would be redundant.

- Any system can be transformed in a Bayesian Network (called a grounded network).
- The generation algorithm is polynomial [Pfeffer, 2000].
- It is possible to use Bayesian Networks state-of-the-art inference algorithms.

- It is easy to produce very large scale systems :
  - naive grounded networks use too much memory;
  - intelligent grounded networks would be redundant.
- Faster inference!



























Relational skeleton

• Each class in a PRM defines a pattern repeated in each of its instances.

- Each class in a PRM defines a pattern repeated in each of its instances.
- In each class, some attributes can be eliminated before others.

- Each class in a PRM defines a pattern repeated in each of its instances.
- In each class, some attributes can be eliminated before others.
- Eliminating such attributes creates an identical potential in each instance.

- Each class in a PRM defines a pattern repeated in each of its instances.
- In each class, some attributes can be eliminated before others.
- Eliminating such attributes creates an identical potential in each instance.
- SVE uses this principle to prevent redundant computations.

- Each class in a PRM defines a pattern repeated in each of its instances.
- In each class, some attributes can be eliminated before others.
- Eliminating such attributes creates an identical potential in each instance.
- SVE uses this principle to prevent redundant computations.
- It does not require the generation of a grounded network.

- Each class in a PRM defines a pattern repeated in each of its instances.
- In each class, some attributes can be eliminated before others.
- Eliminating such attributes creates an identical potential in each instance.
- SVE uses this principle to prevent redundant computations.
- It does not require the generation of a grounded network.
- However, it works only if the inner attributes are not observed.

- Each class in a PRM defines a pattern repeated in each of its instances.
- In each class, some attributes can be eliminated before others.
- Eliminating such attributes creates an identical potential in each instance.
- SVE uses this principle to prevent redundant computations.
- It does not require the generation of a grounded network.
- However, it works only if the inner attributes are not observed.

Once inner attributes are eliminated, SVE proceeds with a **bottom-up** elimination of each instance.















#### Probability P(C)?



Elimination at class level : performed only once !



It works only if the inner attributes are not observed. Only internal variables can be eliminated at class level !



 $\ensuremath{\operatorname{FIGURE}}$  : Different configurations due to evidence.

### Integrating D-Separation in SVE

- Recursive calls are made on each reverse slot chain.
- If at least one reverse slot chain is active :
  - Activate any required slot chains.
  - Either compute the inner node elimination or retrieve any existing computation.



 $\ensuremath{\operatorname{Figure}}$  : Different configurations due to evidence.

### Integrating D-Separation in SVE

- Recursive calls are made on each reverse slot chain.
- If at least one reverse slot chain is active :
  - Activate any required slot chains.
  - Either compute the inner node elimination or retrieve any existing computation.



 $\ensuremath{\operatorname{FIGURE}}$  : Different configurations due to evidence.

### Integrating D-Separation in SVE

- Recursive calls are made on each reverse slot chain.
- If at least one reverse slot chain is active :
  - Activate any required slot chains.
  - Either compute the inner node elimination or retrieve any existing computation.



 $\ensuremath{\operatorname{FIGURE}}$  : Different configurations due to evidence.

### Integrating D-Separation in SVE

- Recursive calls are made on each reverse slot chain.
- If at least one reverse slot chain is active :
  - Activate any required slot chains.
  - Either compute the inner node elimination or retrieve any existing computation.



 $\ensuremath{\operatorname{Figure}}$  : Different configurations due to evidence.

### Integrating D-Separation in SVE

- Recursive calls are made on each reverse slot chain.
- If at least one reverse slot chain is active :
  - Activate any required slot chains.
  - Either compute the inner node elimination or retrieve any existing computation.
# D-Separation in PRMs (monomorphic and)



 $\ensuremath{\operatorname{FIGURE}}$  : Different configurations due to evidence.

### Integrating D-Separation in SVE

Following SVE bottom-up elimination order :

- Recursive calls are made on each reverse slot chain.
- If at least one reverse slot chain is active :
  - Activate any required slot chains.
  - Either compute the inner node elimination or retrieve any existing computation.

# D-Separation in PRMs (monomorphic and)



 $\ensuremath{\operatorname{Figure}}$  : Different configurations due to evidence.

### Integrating D-Separation in SVE

Following SVE bottom-up elimination order :

- Recursive calls are made on each reverse slot chain.
- If at least one reverse slot chain is active :
  - Activate any required slot chains.
  - Either compute the inner node elimination or retrieve any existing computation.

#### Generalized BayesBall rule

#### The example used for our experimentations

- We used the printer example with the following parameters :
  - different number of computers per room;
  - different number of printers per room;
  - different number of evidences.
- The examples we used have a 1400 to 24000 attributes.

## Efficiency of D-Separation



### Increasing the number of computers



## Increasing the number of printers



 $\bullet$  Construction by experts  $\Longrightarrow$  genericity and design patterns.

- $\bullet$  Construction by experts  $\Longrightarrow$  genericity and design patterns.
- $\Longrightarrow$  For instance : in genealogy, relations only between parents and children.

 $\bullet$  Construction by experts  $\Longrightarrow$  genericity and design patterns.

 $\Longrightarrow$  For instance : in genealogy, relations only between parents and children.



 $\bullet$  Construction by experts  $\Longrightarrow$  genericity and design patterns.

 $\Longrightarrow$  For instance : in genealogy, relations only between parents and children.



 $\bullet$  Construction by experts  $\Longrightarrow$  genericity and design patterns.

 $\Longrightarrow$  For instance : in genealogy, relations only between parents and children.



• Solution : create compound classes : dynamic classes

 $\bullet$  Construction by experts  $\Longrightarrow$  genericity and design patterns.

 $\Longrightarrow$  For instance : in genealogy, relations only between parents and children.



• Solution : create compound classes : dynamic classes



## Substitution rule



#### Rule

One random variable can only belong to one instance.

## Substitution rule



#### Rule

One random variable can only belong to one instance.

## Substitution rule



#### Rule

One random variable can only belong to one instance.

#### Proposition

The following problem is NP-hard :

**Instance :** a PRM, an integer  $K \ge 0$ 

**Question :** does there exist a set of dynamic classes / substitutions s.t. the number of operations performed by structured inference is  $\leq K$ ?

#### Proposition

The following problem is NP-hard :

**Instance :** a PRM, an integer  $K \ge 0$ 

**Question :** does there exist a set of dynamic classes / substitutions s.t. the number of operations performed by structured inference is  $\leq K$ ?

remains NP-hard even when classical inference is polynomial

## Approximate algorithm

### Boundary graph





#### Relational skeleton

## Approximate algorithm

### Boundary graph



#### $\Rightarrow$ find frequent patterns in the boundary graph

Modèles Probabilistes Complexes

• Mining frequent subgraphs

[Inokuchi et al, 05], [Kuramochi and Karypis, 01], [Yan and Han, 02]

### Search tree

Mining frequent subgraphs

[Inokuchi et al, 05], [Kuramochi and Karypis, 01], [Yan and Han, 02]

• Variant of the gSpan algorithm :

1 edge 2 edges 3 edges 4 edges



• Node = (dynamic class, set of instances)

### Search tree

Mining frequent subgraphs

[Inokuchi et al, 05], [Kuramochi and Karypis, 01], [Yan and Han, 02]

• Variant of the gSpan algorithm :



• Node = (dynamic class, set of instances)

Mining frequent subgraphs

[Inokuchi et al, 05], [Kuramochi and Karypis, 01], [Yan and Han, 02]

• Variant of the gSpan algorithm :



• Node = (dynamic class, set of instances)

• Possible substitutions (Rule 1) : Max Independence Set

### Efficiency Pruning Rule

 $\label{eq:prune nodes} Prune \ nodes/classes \ whose \ substitutions \ do \ not \ speed-up \ inference.$ 

#### Efficiency Pruning Rule

 $\label{eq:prune nodes} Prune \ nodes/classes \ whose \ substitutions \ do \ not \ speed-up \ inference.$ 

• Gain estimation by dynamic programming

 $\implies \alpha$ -value :  $\alpha > 0 \iff$  class unattractive

search tree not monotonically  $\alpha$ -decreasing!

 ${\bullet}\, {\sf Rule} \ {\sf applied} \ : \ {\sf prune} \ {\sf subtree} \ {\sf whenever} \ \alpha > 0$ 

